

# Unit- and Sequence Test Generation with HOL-TestGen

**Tests et Methodes Formelles**

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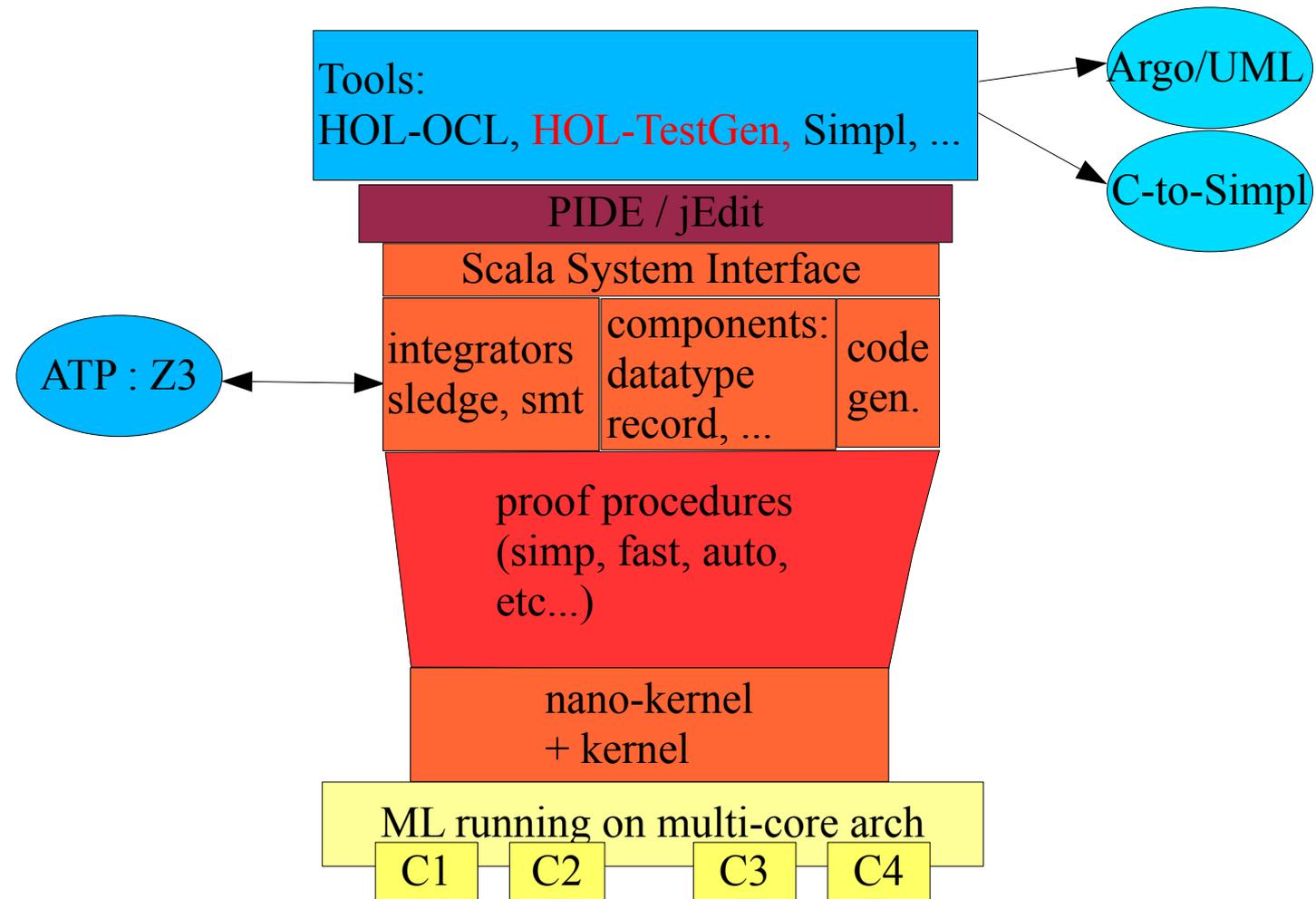
# Overview

- HOL-TestGen and its Business-Case
- The Standard Workflow for Unit Testing
- Demo
- The Workflow for Sequence Tests

# HOL-TestGen and its Business-Case

- HOL-TestGen is somewhat unusual test-Tool:
  - implemented as “PlugIn” in a major Interactive Theorem Proving Environment : Isabelle/HOL
  - conceived as formal testcase-generation method based on symbolic execution of a model (in HOL)
  - Favors Expressivity and emphasizes Test-Plans as formal entities; emphasis on Interactivity

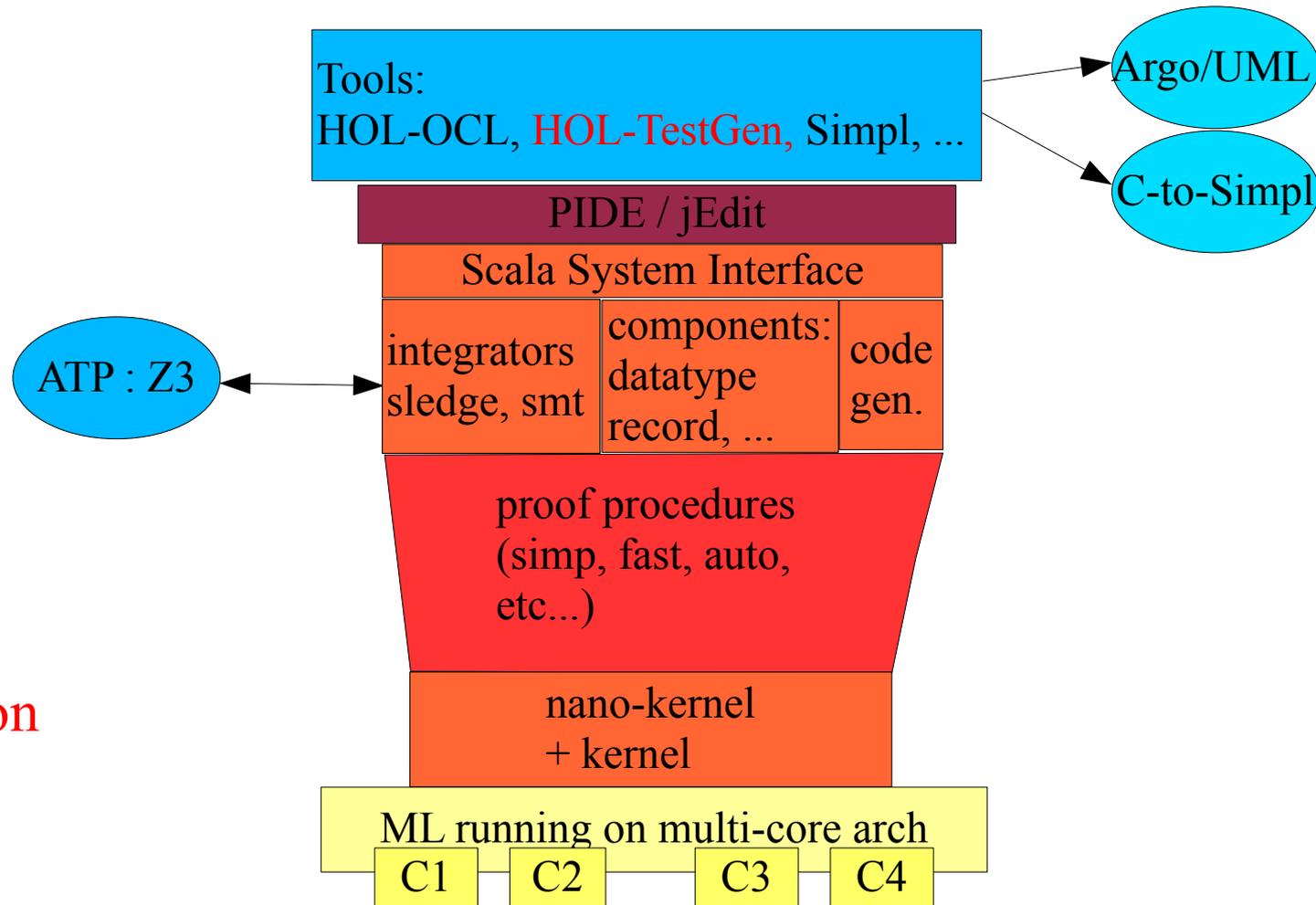
# HOL-TestGen as Plugin in the Isabelle Architecture



# HOL-TestGen as Plugin in the Isabelle Architecture

## Advantage:

- **Reuse** of powerful components in unique, interactive integrated environment
- **seamless integration** of test and proof activities

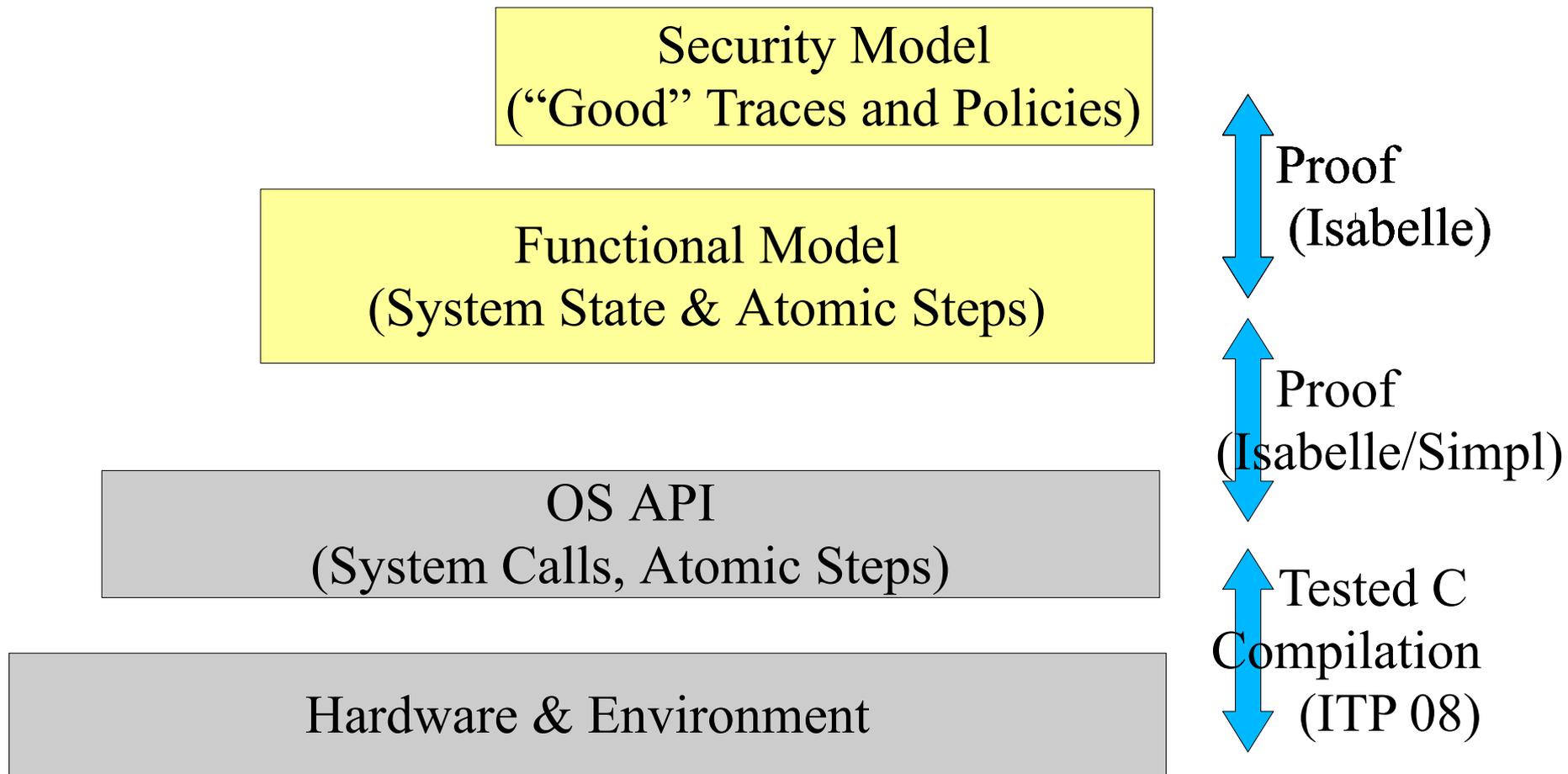


# HOL-TestGen's Business Case

- If you have already a system model in Isabelle or Coq, you might want to link it to a real implementation.
- This has particular relevance in a Certification project (for example: Common Criteria EAL 5 - 7)

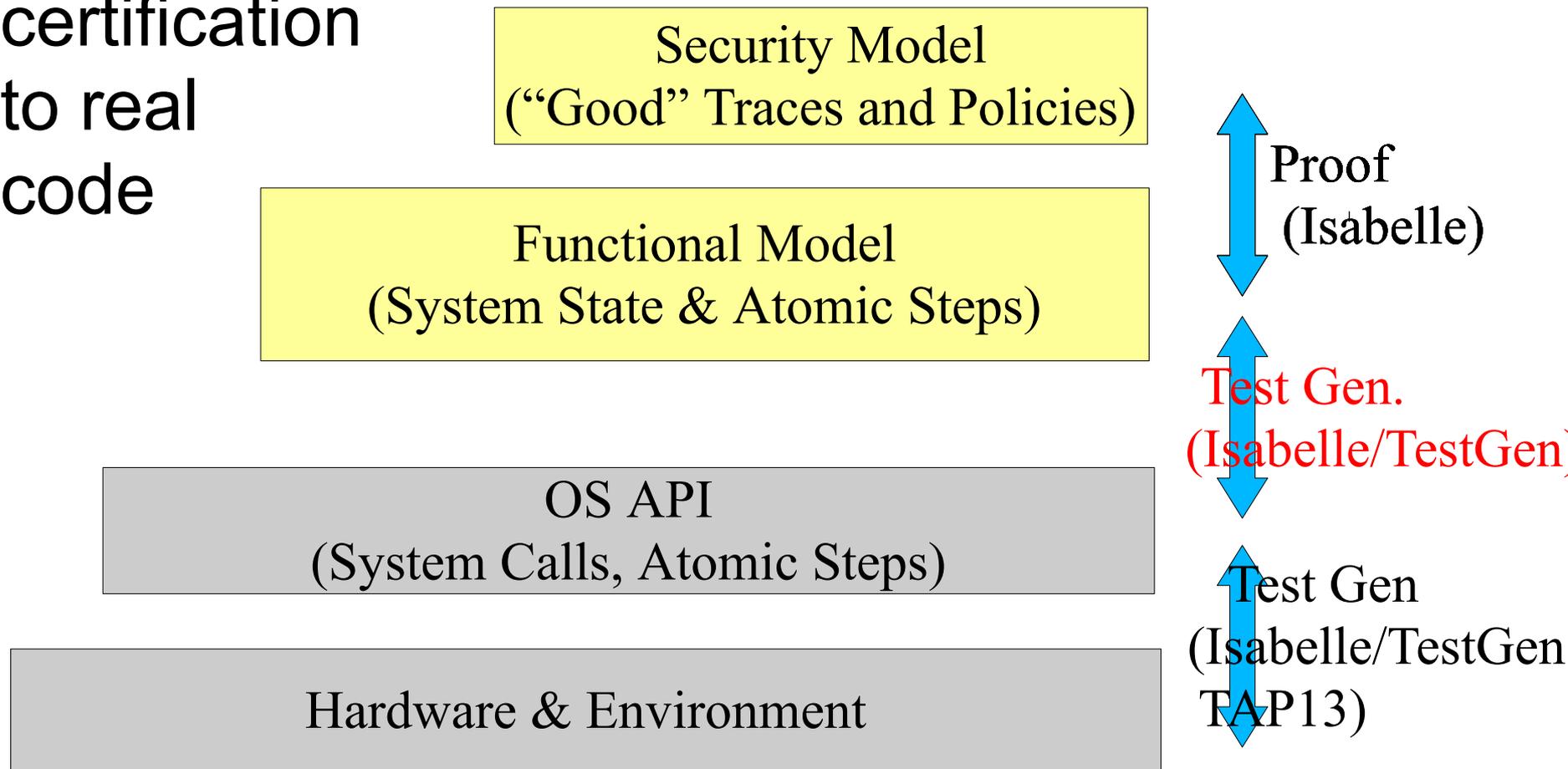
# HOL-TestGen's Business Case

- NICTA seL4 Verified Project

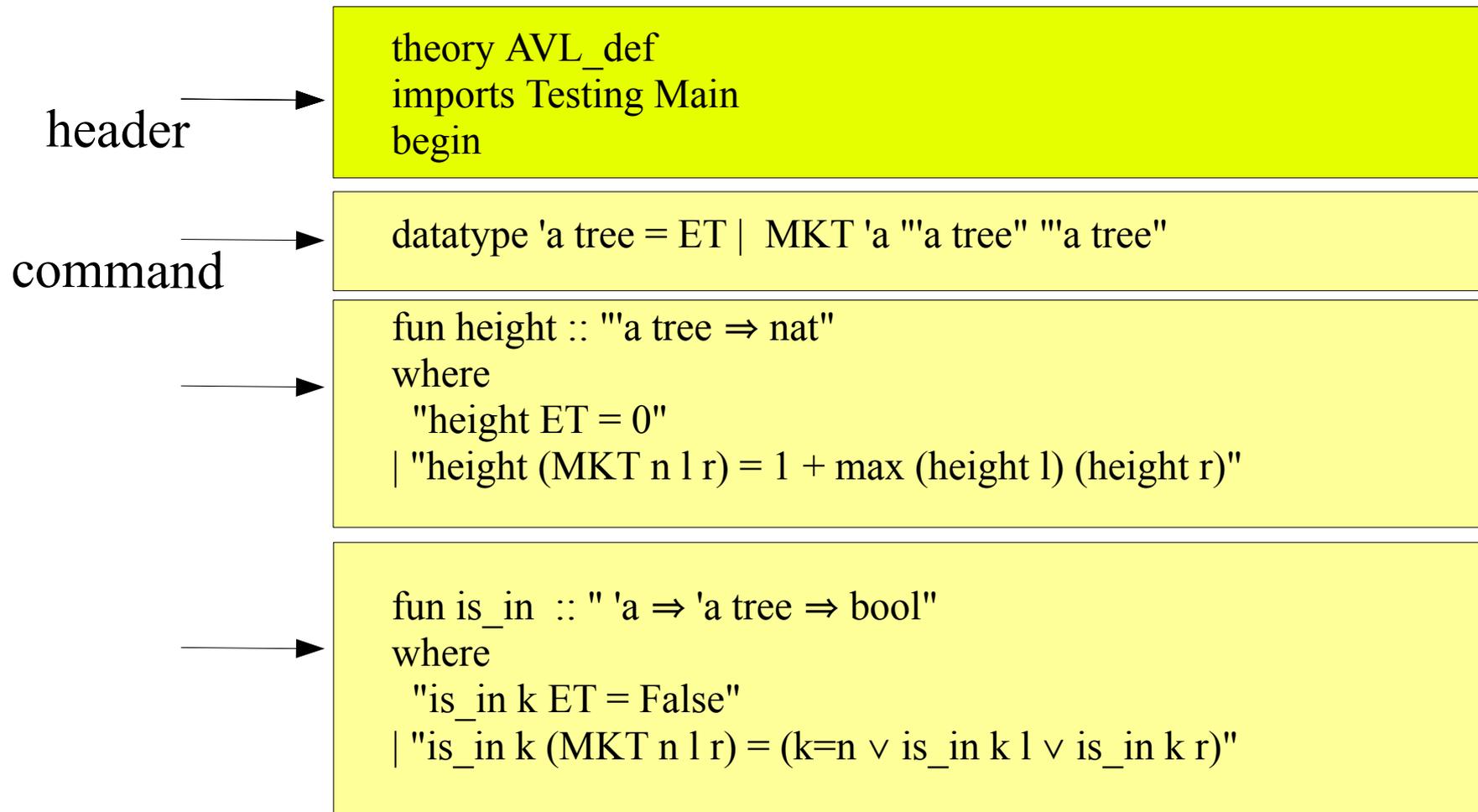


# EUROMILS PikeOS Project

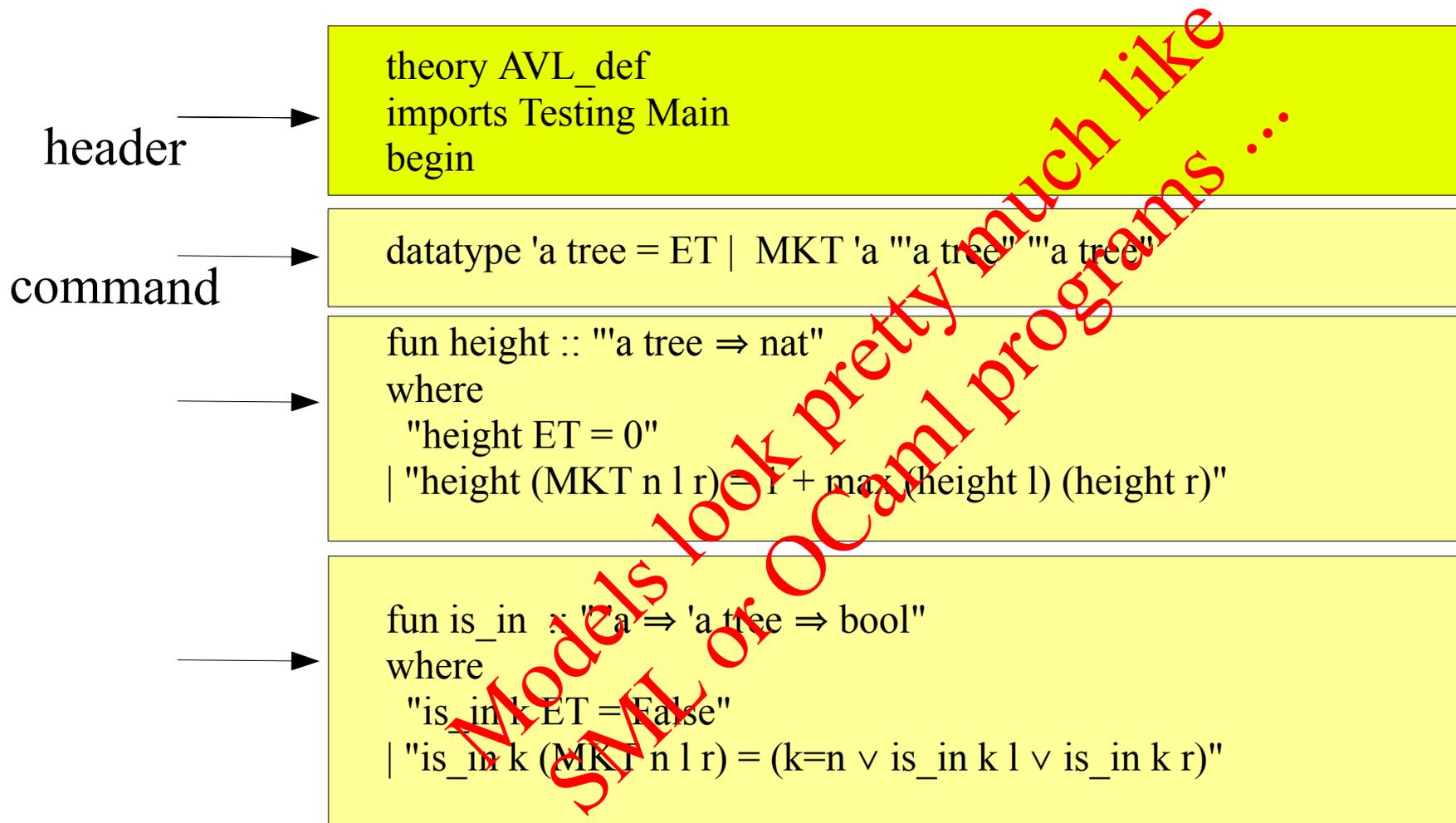
- CC EAL 5+; Linking the FM model of the certification to real code



# Look and Feel : The PIDE Interface for Theories



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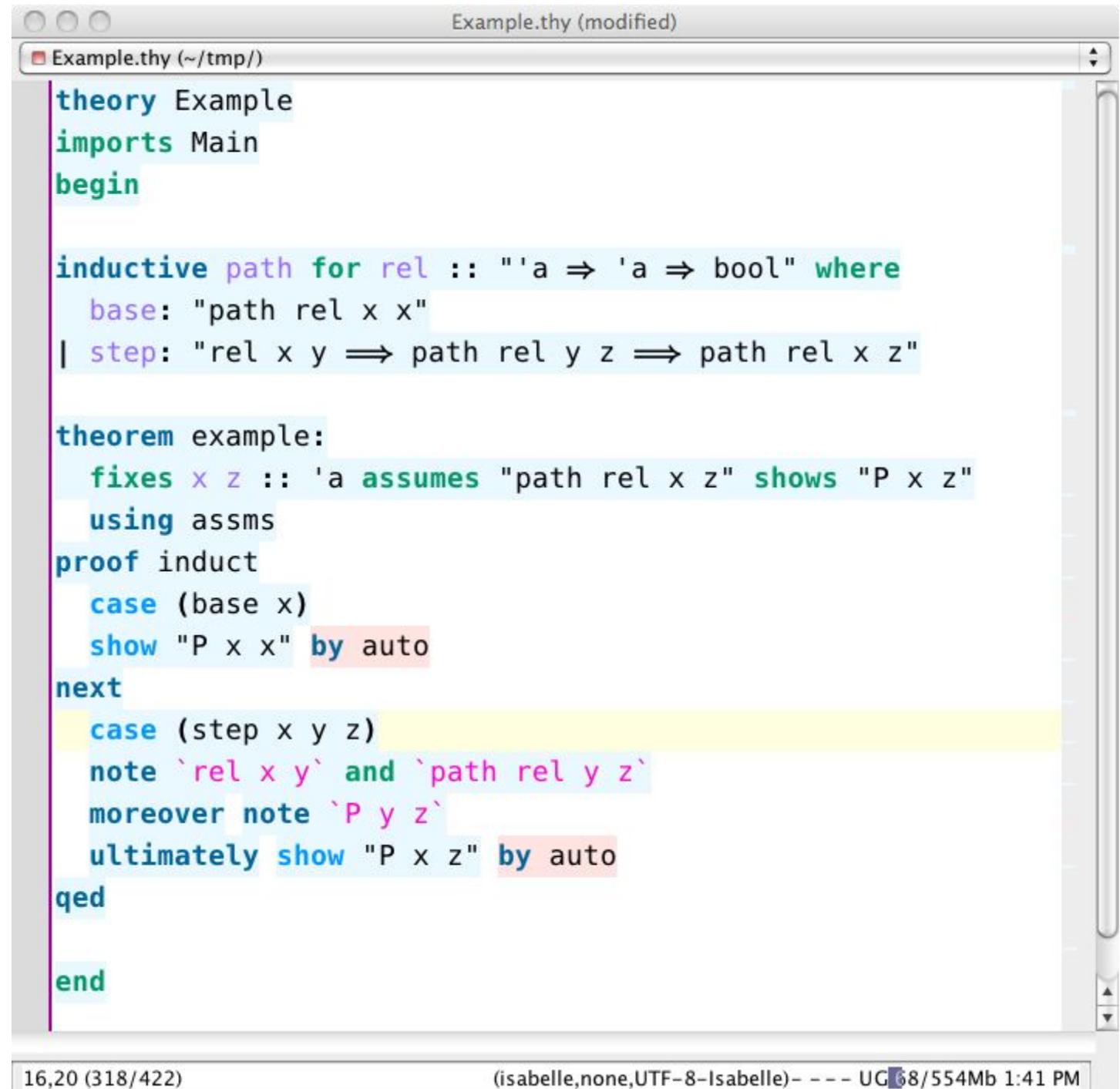


Parallel,  
Asynchronous  
execution and  
validation  
in the

## jEdit - GUI

IDE look and feel;  
Very attractive  
work environment.

(Better than Eclipse ;-)



```
Example.thy (modified)
Example.thy (~/.tmp/)

theory Example
imports Main
begin

inductive path for rel :: "'a ⇒ 'a ⇒ bool" where
  base: "path rel x x"
| step: "rel x y ⇒ path rel y z ⇒ path rel x z"

theorem example:
  fixes x z :: 'a assumes "path rel x z" shows "P x z"
  using assms
proof induct
  case (base x)
  show "P x x" by auto
next
  case (step x y z)
  note `rel x y` and `path rel y z`
  moreover note `P y z`
  ultimately show "P x z" by auto
qed

end
```

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# HOL-TestGen Workflow

- Modelisation
  - writing background theory of problem domain

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory** (the “model”)

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- Writing a **test-theory** (the “model”)

Example: Sorting in HOL

```
primrec is_sorted :: "int list  $\Rightarrow$  bool"  
where  "is_sorted [] = True"  
      | "is_sorted (x#xs) =  
        case xs of  
          []  $\Rightarrow$  True  
        | (y#ys)  $\Rightarrow$  (x  $\leq$  y)  $\wedge$  is_sorted ys"
```

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory** (the “model”)

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**testspec** “ $\text{is\_sorted}(PUT\ x)$   
 $\wedge \text{asc}(x, PUT\ x)$ ”

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification** TS

**pattern:**

**testspec** “pre  $x \rightarrow$  post  $x$  ( $PUT\ x$ )”

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**

**example:**

`test_spec “is_sorted  $x$   $\rightarrow$  is_sorted (prog a x)”`

or

`test_spec “is_sorted (PUT l)”`

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**  
(“Testcase Generation”)

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```
apply(gen_test_cases 3 1 “PUT”)
```

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**

(“Testcase Generation”)

$$TC_1 \Rightarrow \dots \Rightarrow TC_n \Rightarrow \text{THYP}(H_1) \Rightarrow \dots \Rightarrow \text{THYP}(H_m) \Rightarrow \text{TS}$$

- where testcases  $TC_i$  have the form

$$\text{Constraint}_1(x) \Rightarrow \dots \Rightarrow \text{Constraint}_k(x) \Rightarrow P(\text{prog } x)$$

- and where  $\text{THYP}(H_i)$  are test-hypothesis

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**

Example:

is\_sorted (PUT I)

1: is\_sorted(PUT [])

2: is\_sorted(PUT [?X])

3: THYP( $\exists x. \text{is\_sorted}(\text{PUT } [x]) \rightarrow \forall x. \text{is\_sorted}(\text{PUT } [x])$ )

4: is\_sorted(PUT [?X, ?Y])

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**

...

5:  $\text{THYP}(\exists x y. \text{is\_sorted}(\text{PUT}[x,y]) \rightarrow \forall x y. \text{is\_sorted}(\text{PUT}[x,y]))$

6:  $\text{is\_sorted}(\text{PUT} [?X, ?Y, ?X])$

7:  $\text{THYP}(\exists x y z. \text{is\_sorted}(\text{PUT} [x,y,z]) \rightarrow \forall x y z. \text{is\_sorted}(\text{PUT} [x,y,z]))$

8:  $\text{THYP}(3 < || \rightarrow \text{is\_sorted}(\text{PUT } I))$

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
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- Conversion into **test-theorem**
- Generation of **test-data**

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- Writing a **test-theory**
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- Generation of **test-data**

`gen_test_data “...”`

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**
- Generation of **test-data**

```
is_sorted(PUT 1 [])  
is_sorted(PUT 1 [0])  
is_sorted(PUT 1 [2])  
is_sorted(PUT 1 [1,2])
```

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**
- Generation of **test-data**
- Generating a **test-harness**

# Black-Box Testing: “The Standard Workflow”

- Writing a **test-theory**
- Writing a **test-specification TS**
- Conversion into **test-theorem**
- Generation of **test-data**
- Generating a **test-harness**
- Run of testharness and  
generation of **test-document**

# Midi Example: Red Black Trees

## Red-Black-Trees: Test Specification

```
testspec :  
(redinv t ^  
 blackinv t)
```

→

```
(redinv (delete x t) ^  
 blackinv (delete x t))
```

where `delete` is the program under test.

# HOL-TestGen Workflow

Demo

# Black-Box Sequence Testing:

- HOL is a state-less language;  
how to model and test stateful systems ?
- How to test systems where you have only control over the initial state ?
- How to test concurrent programs implementing a model ?

# How to model and test stateful systems ?

- Use Monads !!!

- The transition in an automaton  $(\sigma, (\iota, o), \sigma)$  set can isomorphically be represented by:

$$\iota \Rightarrow \sigma \Rightarrow (o, \sigma) \text{ set}$$

or for a deterministic transition function:

$$\iota \Rightarrow \sigma \Rightarrow (o, \sigma) \text{ option}$$

... which category theorists or functional programmers would recognize as a **Monad function space**

# How to model and test stateful systems ?

- Use Monads !!!
  - The transition in an automaton  $(\sigma, (l, o), \sigma)$  set can isomorphically be represented by:

$$l \Rightarrow (o \times \sigma) \text{ Mon}_{\text{SBE}}$$

or for a deterministic transition function:

$$l \Rightarrow (o \times \sigma) \text{ Mon}_{\text{SE}}$$

... which category theorists or functional programmers  
would recognize as a **Monad function space**

# How to model and test stateful systems ?

- Monads must have two combination operations bind and unit enjoying three algebraic laws.
  - For the concrete case of  $\text{Mon}_{\text{SE}}$ :

```

definition bindSE :: "('o, 'σ)MONSE ⇒ ('o ⇒ ('o', 'σ)MONSE) ⇒ ('o', 'σ)MONSE"
where      "bindSE f g = (λσ. case f σ of None ⇒ None
                                     | Some (out, σ') ⇒ g out σ')"
  
```

```

definition unitSE :: "'o ⇒ ('o, 'σ)MONSE" ("(return _)" 8)
where      "unitSE e = (λσ. Some(e, σ))"
  
```

– and write  $o \leftarrow m; m' o$  for  $\text{bind}_{\text{SE}} m (\lambda o. m' o)$   
 and  $\text{return}$  for  $\text{unit}_{\text{SE}}$

# How to model and test stateful systems ?

- Valid Test Sequences:

$$\sigma \models o_1 \leftarrow m_1 \iota_1; \dots; o_n \leftarrow m_n \iota_n; \text{return}(P \ o_1 \cdots o_n)$$

- ... can be generated to code
- ... can be symbolically executed ...

$$\frac{}{(\sigma \models \text{return } P) = P}$$

$$\frac{C_m \iota \sigma \quad m \iota \sigma = \text{None}}{(\sigma \models ((s \leftarrow m \iota; m' s))) = \text{False}}$$

$$\frac{C_m \iota \sigma \quad m \iota \sigma = \text{Some}(b, \sigma')}{(\sigma \models s \leftarrow m \iota; m' s) = (\sigma' \leftarrow (m' b))}$$

# How to model and test stateful systems ?

- Test Refinements for a step-function SPEC and a step function SUT:

$$\begin{array}{l} \sigma \models o_1 \leftarrow \text{SPEC}_1 \iota_1; \dots; o_n \leftarrow \text{SPEC}_n \iota_n; \text{return}(res = [o_1 \cdots o_n]) \\ \rightarrow \\ \sigma \models o_1 \leftarrow \text{SUT}_1 \iota_1; \dots; o_n \leftarrow \text{SUT}_n \iota_n; \text{return}(res = [o_1 \cdots o_n]) \end{array}$$

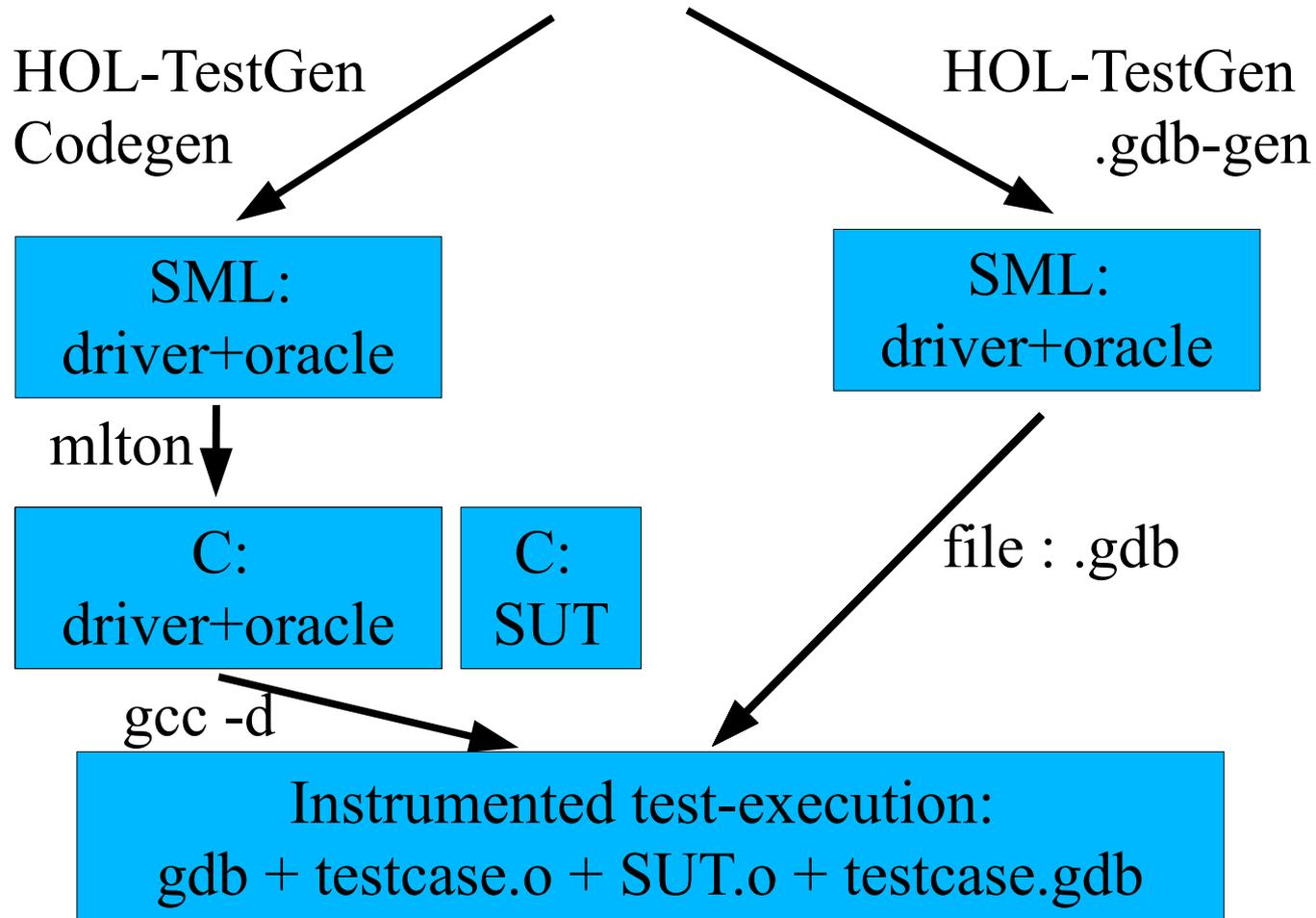
- The premis is reduced by symbolic execution to constraints over  $res$ ; a constraint solver (Z3) produces an instance for  $res$ . The conclusion is compiled to a test-driver/test-oracle linked to  $SUT$ .

# How to test concurrent programs implementing a model ?

- Assumption: Code compiled for LINUX and instrumented for debugging (gcc -d)
- Assumption: No dynamic thread creation (realistic for PikeOS); identifiable atomic actions in the code;
- Assumption: Mapping from abstract atomic actions in the model to code-positions known.
- Abstract execution sequences were generated to .gdb scripts forcing explicit thread-switches of the SUT executed under gdb.

# How to test concurrent programs implementing a model ?

$\sigma \models o_1 \leftarrow \text{SUT}_1 \iota_1; \dots; o_n \leftarrow \text{SUT}_n \iota_n; \text{return}(res = [o_1 \cdots o_n])$



# Conclusion

- HOL-TestGen is an Advanced Model-based Testing Environment built on top of Isabelle/HOL
- Allows to establish a Link between a formal System Model in Isabelle/HOL and Real Code by (semi)-automated generation of tests.
- Smooth Integration of Test and Proof !