Formal Methods at Airbus: Experience Feedback

Presented by Jean Souyris / EYYWDV – Verification and dependability support



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- Introduction
- Context, Objectives and Constraints
- Basics
- Industrial state-of-the-use
- Next transfers
- Main issues: solved & remaining
- Conclusion



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Context: avionics software products

Avionics domains

- Flight controls: safety-critical (DAL A), time-critical, SCADE specification (synchronous paradigm), no operating system, floating-point calculus, and also non SCADE "driver-like" functions
- Flight Warning: medium criticality (DAL C), asynchronous (multitasks) functions running on IMA platform, complex data structures (non dynamic allocation)
- **Board/ground communication**: medium criticality (DAL C), asynchronous (multi-tasks) functions running on IMA or POSIX platforms, complex data structures (no dynamic allocation)
- Maintenance functions: low criticality (DAL D & E), asynchronous (multi-tasks) functions running on POSIX OS

Verification environment

• **SIMUGENE**: hardware virtualization for verification by execution



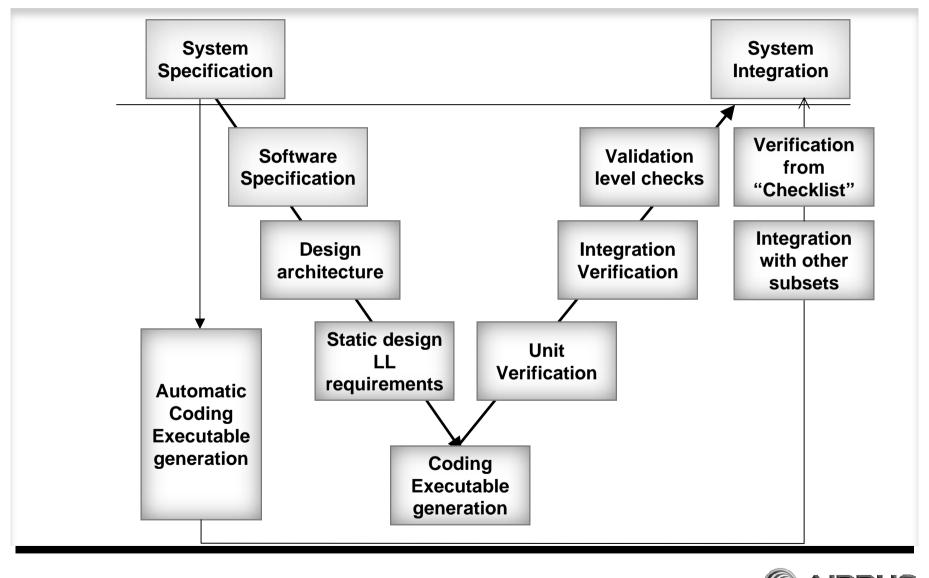
Context: elements of development organisation

Avionics software development teams

- Specify, design, code and verify software products from system specifications
- In conformance with Airbus's reference development processes and methods, thus with DO-178B
- **Support teams** (specification, design, verification, configuration management, modification management)
 - Strategies
 - Operational support Methods and tools (including training)
 - (new) Service activities on behalf of development teams
- Process and assurance teams ("Quality")
 - Process definition
 - Check the conformance with reference process and DO-178B



Avionics software development Process





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Context / Objectives for formal tools

• Steady increase of System complexity

- Master verification costs
- Performance: contribute to the safe and optimal use of modern hardware and software features
- Keep computation safety (executability) verification at high level

• Need for early maturity

- Exhaustive verification techniques
- Available as soon system design / code is available

Long term product durability and maintainability

- Localized modifications and automatic replay
- Postpone hardware re-engineering by optimal resource usage analysis

Towards Calculus Based Engineering and Product Based Assurance



Constraints

- Soundness
- High Automaticity and scalability
 - In intended usage domain at airbus's
- Analysis of unaltered programs
 - "What is analysed is what will fly"
- Usability by standard software engineers on standard machines
 - No initial high level skills in theoretical computer science required
 - Standard workstations Airbus uses to buy
- Ability to be integrated into the DO178B (and C !) conforming process



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Basic principles

- Abstract Interpretation based Static Analysis
 - <u>http://www.di.ens.fr/~cousot/Al/</u>
- Program Proof
 - Hoare's triple: <u>http://en.wikipedia.org/wiki/Hoare_logic</u>
 - Dijkstra's Weakest Precondition: <u>http://en.wikipedia.org/wiki/Predicate_transformer_semantics</u>
 - Automatic theorem proving: <u>http://ergo.lri.fr/</u>



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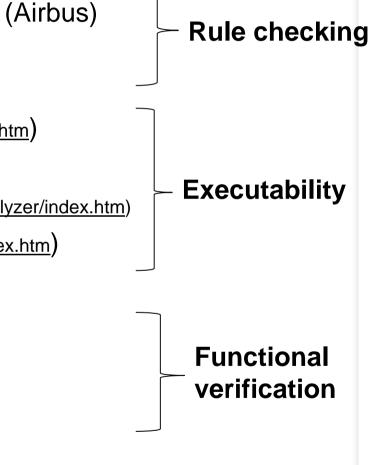
Verification Tools

Static Analysis

- Frama-c (CEA, <u>http://frama-c.com/</u>) / TASTER (Airbus)
- Frama-c / FAN-C (Airbus)
- Astrée (AbsInt, http://www.absint.com/astree/index.htm)
- Fluctuat (CEA)
- a3 / Stack(AbsInt, http://www.absint.com/stackanalyzer/index.htm)
- a3 / WCET (AbsInt, http://www.absint.com/ait/index.htm)

• Program Proof (deductive methods)

• Caveat (CEA)





Rule checking

CheckC / TASTER

- Functionality: Verification of C coding rules
- On top of Frama-c kernel (exploits the AST built by the kernel)

• Fan-C

- Functionality: Verification of control and data flows (conformity LLR <-> C code)
- Abstract Interpretation based static analysis of the C source
- Takes profit from Frama-C Kernel (AST CIL) and plug-ins : Value, Users, Inout and From



Executability

- Astrée (AbsInt, ENS <u>http://www.astree.ens.fr/</u>)
 - Functionality: proof of absence of Run Time Errors of C programs
 - Abstract Interpretation based static analysis of the C source code
 - "Double specialisation" paradigm for precision ("zero false alarm")
 - Best suited for embedded synchronous C programs produced from "SCADE like" specifications

• Fluctuat (CEA)

- Functionality: computes floating-point inaccuracies, proves stability computation schemes, performs some functional proofs
- Abstract Interpretation based static analysis of the C source code
- Best suited for the analysis library components



Executability

- a³ / Stack (<u>http://www.absint.com/ait/index.htm</u>)
 - Functionality: computes an upper-bound of the memory consumed by the program stack (usually from a task's entry point)
 - Maxim memory allocated to the stack is set accordingly
 - Static analysis by Abstract Interpretation of programs in binary form
- a³ / WCET (<u>http://www.absint.com/stackanalyzer/index.htm</u>)
 - Functionality: computes an upper-bound of the Worst Case Execution Time (usually from a task's entry point)
 - This upper-bound can then be compared to an allowed time-budget
 - Static analysis by Abstract Interpretation of programs in binary form
 - Includes a model of the processor and peripherals
 - Best suited for embedded synchronous C programs produced from "SCADE like" specifications



Program proof

• Caveat (CEA)

- Functionality: Proof of specifications expressed in first order logic
- Analysis of C source code
- Weakest Precondition (Dijkstra) computation
- Theorem proving (Caveat's theorem prover + Alt-Ergo (INRIA))
- Best suited for source code vs Low Level requirements verification



Current scope of the tools (+method)

	Flight Controls (DAL A)	(Platform) DAL B & C functions	(Platform) DAL D functions	Platform software (drivers)	I/O boards	DAL E or Software tools
Rule checking	~	\checkmark	~	~	~	~
Executability	~	√12	√12	√12	~	√12
Program proof	~	~	~	~	~	~

¹: RTE (Astrée, CodeSonar), Floating-point (Fluctuat)

²: Stack usage (a3 / Stack)

³: a3 / WCET

Tools (+method of use) vs objectives

	Perfo ¹	Computation safety	No other (sound) mean	Activity cost savings	Early maturity	Product durability
Rule checking	-	-	-	~	✓	-
Executability	√ 1	\checkmark	~	-	~	√1
Program proof	-	~	-	~	~	-

¹: contribution to the optimal use of hardware resources: a³ / Stack and WCET



Tools (+ methods) vs constraints

	Soundness	Automaticity & scalability	Unaltered programs	Standard engineers	Standard machines	DO-178
Rule checking	√1	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
Executability	√1	~	√2	√3	~	√4
Program proof	~	~	~	~	~	✓

- ¹: With the exception of syntactic and pattern matching tools
- ²: Some pieces of code like asm blocks must be removed (rare); insertion of directives
- ³: Astrée, Fluctuat: service currently performed by static analysis specialists ;
- ⁴: So far, the decision to claim a certification credit from the use of Astrée and Fluctuat has not been made;



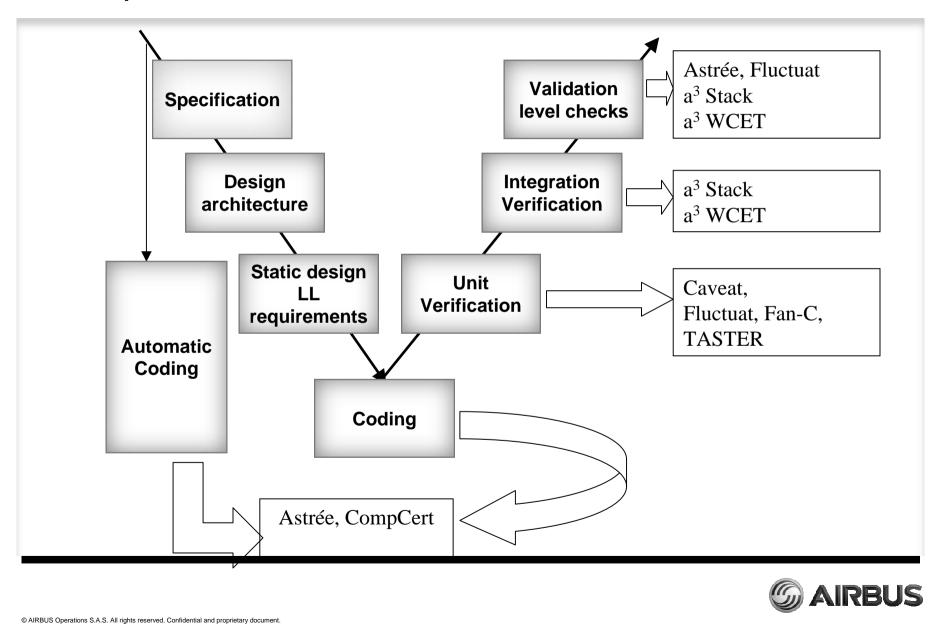
Current Deployment

	Flight Controls (DAL A)	(Platform) DAL B&C functions	(Platform) DAL D functions	Platform software (drivers)	I/O boards	DAL E or Software tools
Rule checking	√1	√ 12	√12	√1	√1	-
Executability	√345	√34	√34	√4	√ 45	-
Program proof	~	-	-	-	-	~

- ¹: Coding rule checker (CheckC/TASTER)
- ²: Data & Control flow checker (Fan-C)
- ³: Astrée, CodeSonar, Fluctuat
- ⁴ : a3 / Stack
- ⁵ : a3 / WCET



Development Process



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Aerospace Valley Forum on Formal Methods	PowerPC, MacOS X			
CompCert (INRIA, <u>htt</u>	PowerPC, Linux PowerPC, EABI, with GNU or Unix tools PowerPC, EABI, with Diab tools ARM, Linux IA32 (x86 32 bits), Linux IA32 (x86 32 bits), BSD IA32 (x86 32 bits), MacOS X IA32 (x86 32 bits), Cygwin environment under Windows			
 Functionality Optimising C compiler for Targets 				
 Underlying principles & rechnology C compiler developed and proved in Coq 				

• First application domain (EYYW)

- EYYW's interest in CompCert
 - Under control optimisations => WCET reduction
 - Proofs made on source still hold after compilation
- Ricardo Bedin França's CIFRE Thesis (Airbus / IRIT)
- Ongoing feasibility study for application to a flight control function



AstréeA (Ecole normale supérieure, <u>http://www.astreea.ens.fr/</u>)

• Functionality

• Proof of absence of Run Time Errors of asynchronous programs

• Underlying principles & Technology

- Abstract Interpretation based static analysis of the C source code
- Included: a model of the ARINC 653 parallel model

Targeted application domain (EYYW)

- IMA functions (e.g: Flight Warning)
- POSIX functions



Dynamic Analysis

Functional Verification

- Properties expressed formally, i.e., in ACSL (Frama-C specification language)
- Execution on SIMUGENE
- Evaluation of properties rather than proof
- First tool (internal research prototype)
 - Low Level Requirement functional coverage for DAL C function
 - Automation of an heavy intellectual analysis
 - Run time data are captured during execution on SIMUGENE
 - Evaluation is then performed



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Main issues: solved & remaining

• Context

- All formal tools Airbus uses come from research
- Airbus has been working with the researchers and tool developers from the beginning

Solved

- Peculiarities of embedded code (very often low level code)
- Conformance to DO-178B
- Acceptance by developers and managers

Remain to do for benefiting more from Formal Methods

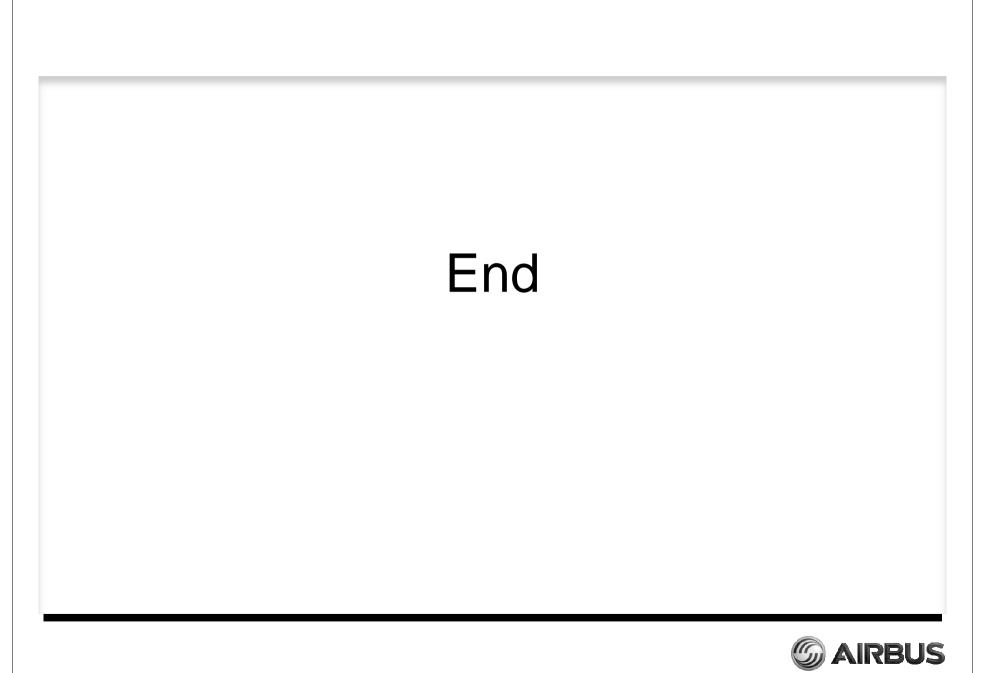
- Proof confirmation after compilation (semantic preservation)
- Deeper process transformation: towards much more computation based engineering



Conclusion

- Ongoing research about a new development strategy
 - Rule checking and executability as soon as code is available
 - Functional verification and coverage by a **combination** of
 - Proof
 - Requires formalised requirements
 - Dynamic analysis
 - The oracles are the formalised requirements
 - Classical test
 - Process definition
 - will still comply with DO-178[BC]...
 - Perhaps without being fully structured by the standard (as it is now)





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