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# Preuve unitaire d'un logiciel avionique

« Preuve de modèle, preuve de programme »

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CYCLE DE CONFÉRENCES TECHNIQUES SUR LES  
MÉTHODES FORMELLES DE DÉVELOPPEMENT

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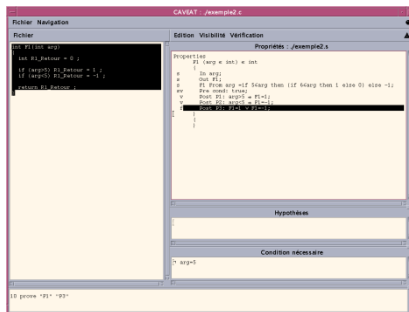
04.02.2014

- ▶ Proof of program overview
- ▶ Unit Proof & Integrated Proof
- ▶ Unit Proof objectives
- ▶ Unit Proof solution
- ▶ Unit Proof results
- ▶ Conclusion

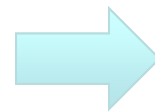
Industrial feedback of Program verification of Airbus avionics software.



Tooling

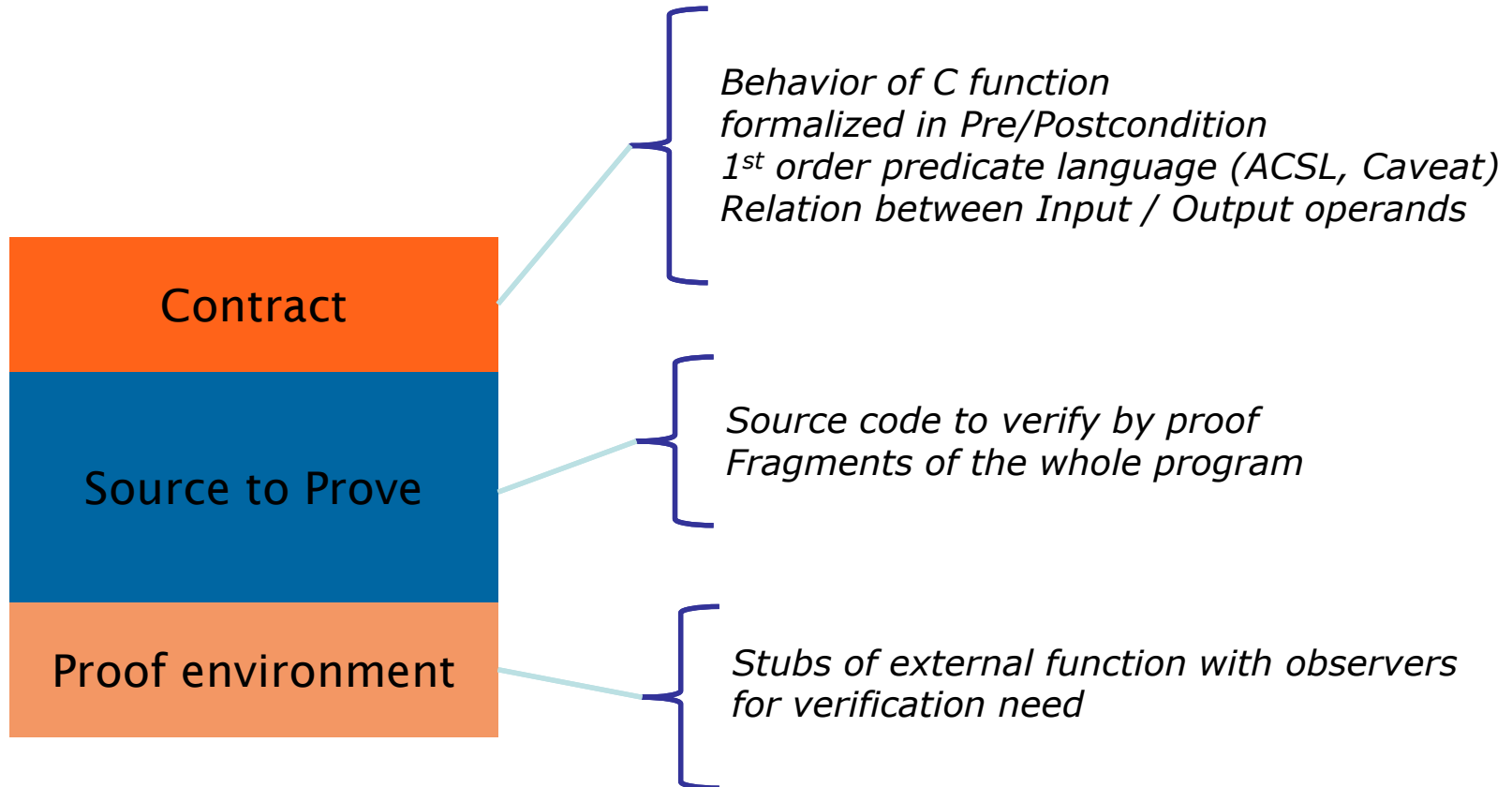


Caveat

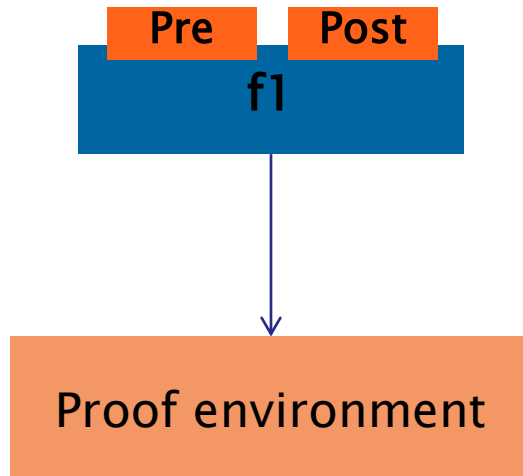


Software Analyzers

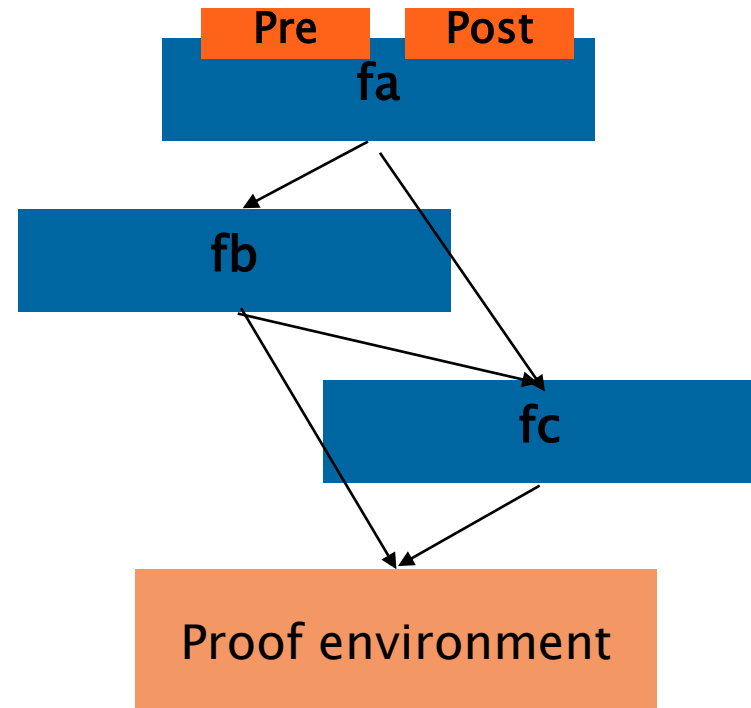
Framac-C/WP



## ▶ Unit Proof



## ▶ Integrated Proof



# Integrated Proof / Example

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```
int MyTab[10] ;
int S;

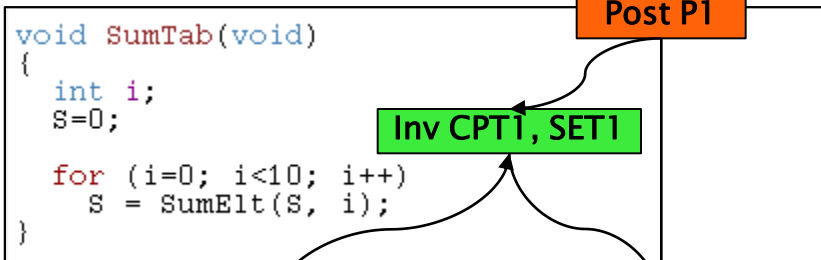
extern int SumElt(int s, int i);
extern int IsEven(int n);
```

```
Const f_even ∈ int -> Bool =
  lambda n ∈ int. (n modulo 2) = 0;

Const f_sum ∈ int -> Tab&int -> int ;

Axiom AX1_sum :
  forall n ∈ int, t ∈ Tab&int.
    n<0 => f_sum(n, t) = 0 ;

Axiom AX2_sum :
  forall n ∈ int, t ∈ Tab&int.
    n>=0 => f_sum(n, t) = f_sum(n-1, t) + (if f_even(t.[.n]) then 1 else 0) ;
```

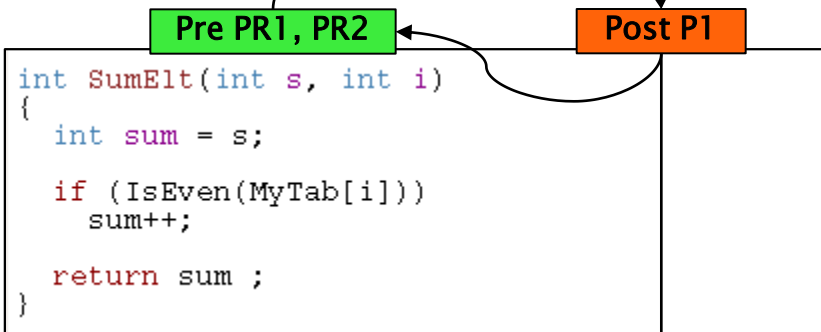


**FONCTION SumTab**

**Post P1** :  $S = f\_sum(9, MyTab)$ ;

**Inv 1 I1\_CPT1** :  $i \geq 0 \ \&\& \ i < 10$  ;

**Inv 1 I1\_SET1** :  $S = f\_sum(i-1, MyTab)$ ;

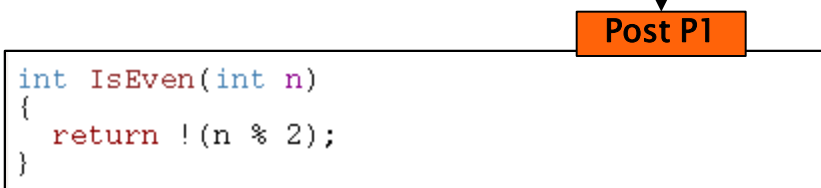


**FONCTION SumElt**

**Pre PR1** :  $s = f\_sum(i-1, MyTab)$ ;

**Pre PR2** :  $i \geq 0$  ;

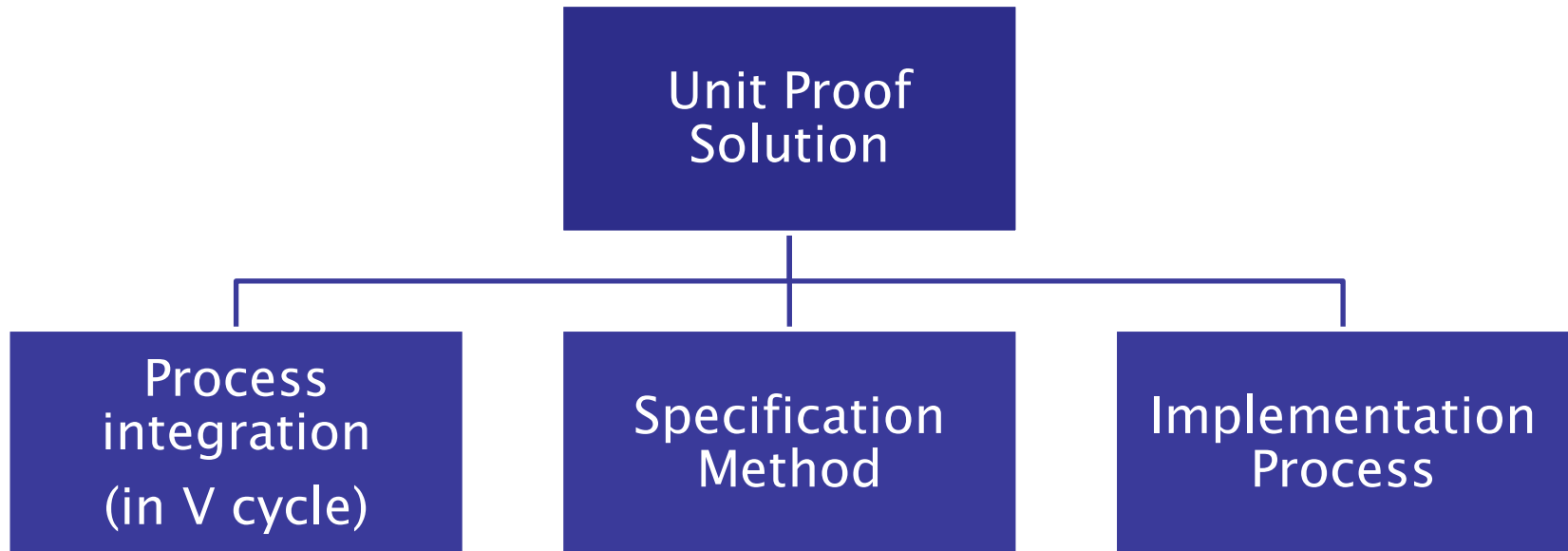
**Post P1** :  $SumElt = f\_sum(i, MyTab)$ ;



**FONCTION IsEven**

**Post P1** :  $IsEven = \text{if } (f\_even(n)) \text{ then } 1 \text{ else } 0$ ;

- ▶ Functional verification of C function (replacing Unit Testing)
- ▶ Achievable for standard developer, non-specialist of formal method
- ▶ Cost Effective
- ▶ Maximal automaticity
- ▶ Quality equivalent (compared with tests)



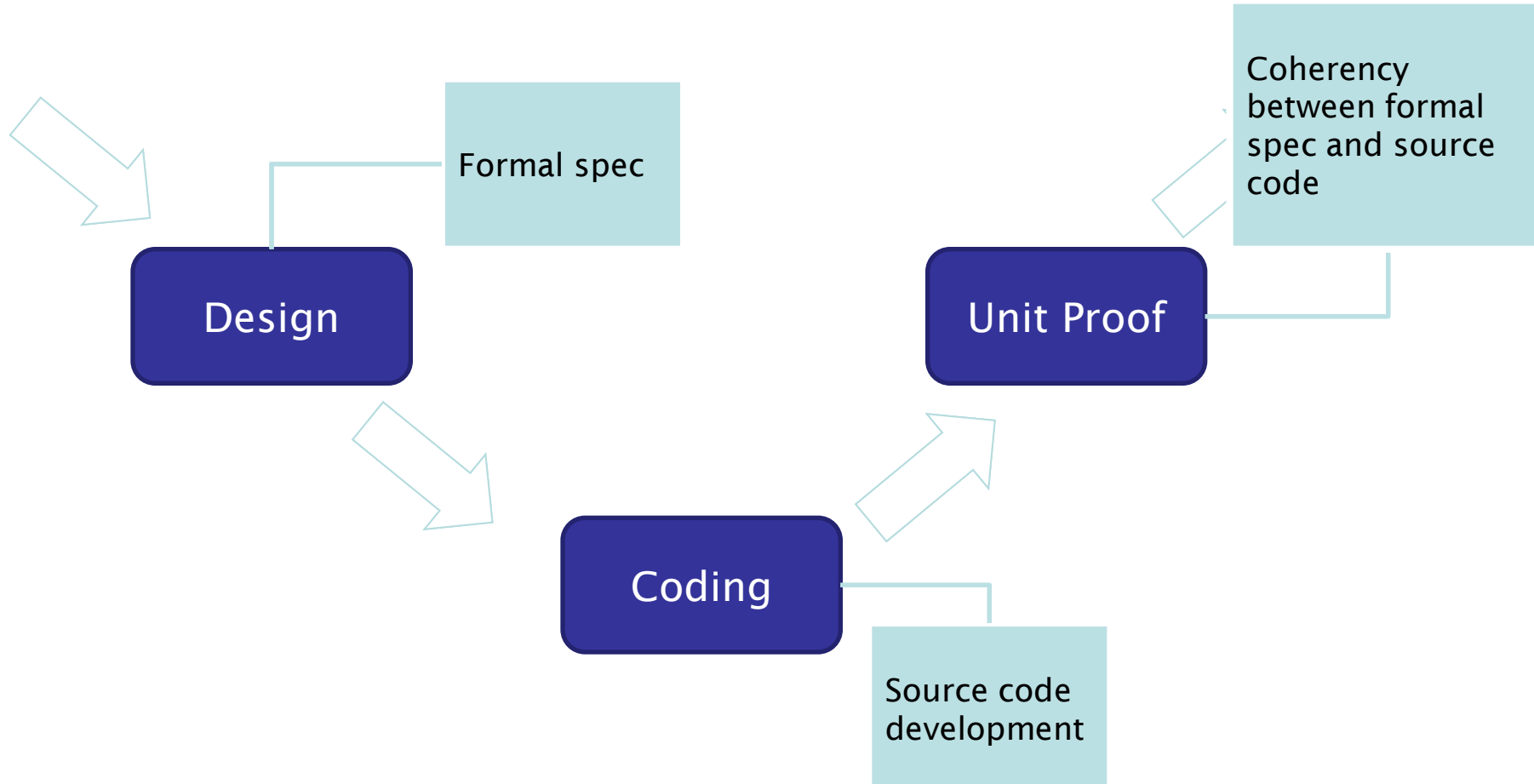


# Unit Proof Solution

## Process Integration

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# Unit Proof Solution

## Specification Method

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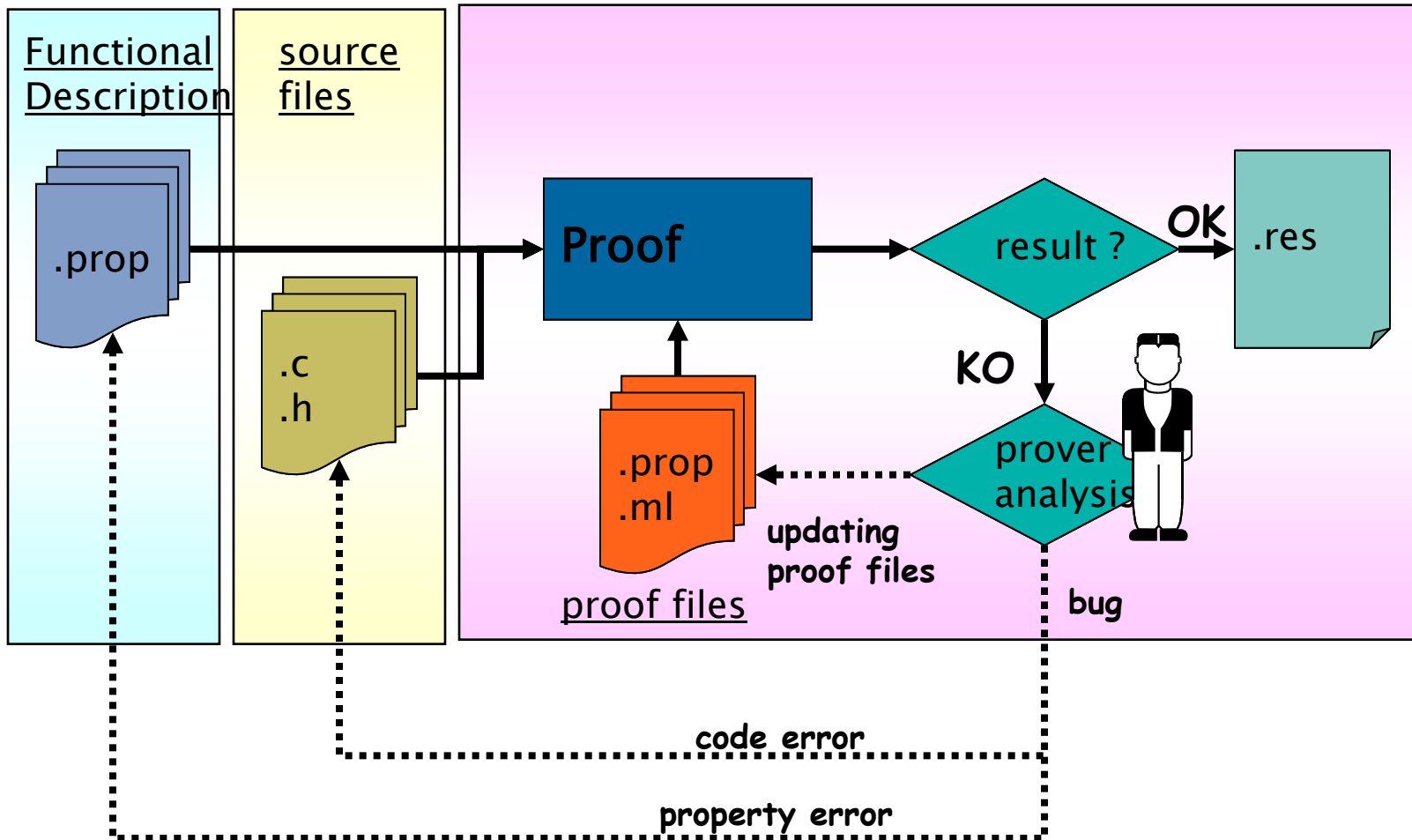
- ▶ Specification of
  - All behaviors
  - All operands
    - Function parameter
    - Global variables
    - In / Out operands of functions called
    - Volatile access
  - Sequence of function calls
  - Constant values
  
- ▶ Definition of a standard way of specification
  - Allowing check activities
  - Easy to read

# Unit Proof Solution

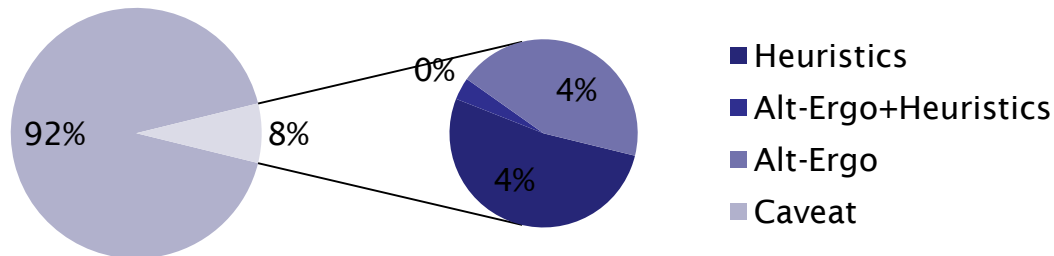
## Implementation Process

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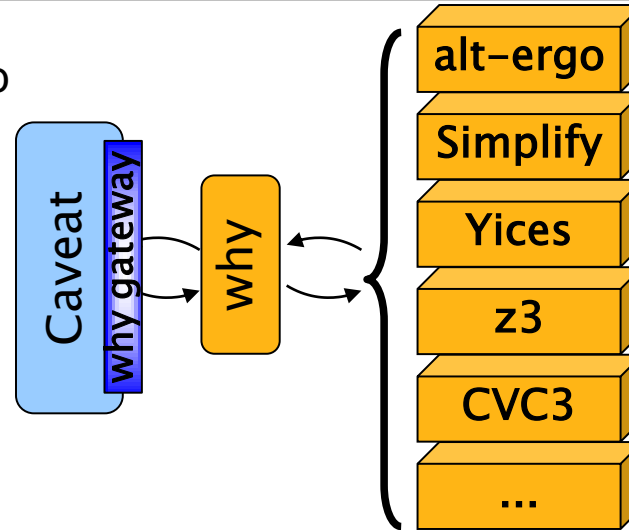


- ▶ Since 2006
- ▶ 3 Aircraft Programs (A380, A400M, A350)
- ▶ ~ 40 kloc of C source Code
- ▶ ~ 5000 Postconditions
- ▶ Proof campaign : ~1 day
- ▶ Automaticity



- ▶ Ease of maintenance
- ▶ No problem of skills in development team

► Caveat Update with AltErgo



► Frama-C NUPW

- Import from Caveat
- Improvement for proof
- Wider area of C programs
- ACSL specification

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## Thank you

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